

No Instruction Computing Using Pointers and Operations in Registers for Adaptable Architecture

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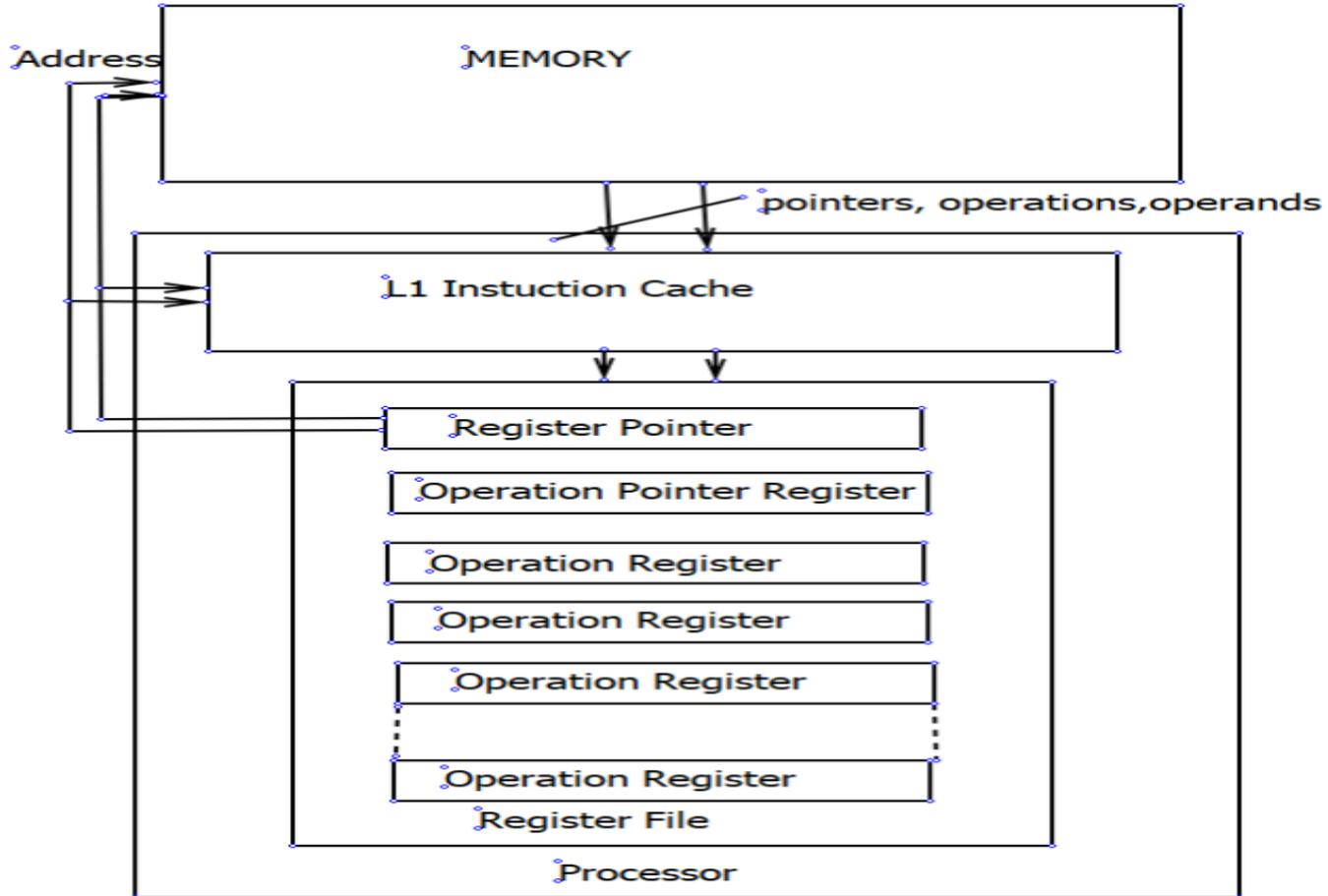
Introduction

- Current architectures depend on delivering maximum numbers of instructions per cycle
- Performance more dependent on memory bandwidth
- Must use program counter each cycle to fetch instructions from memory
- Issuing parallel instructions is complicated
- Operations have more localities than data and could move to registers

No Instructions Computing Concept

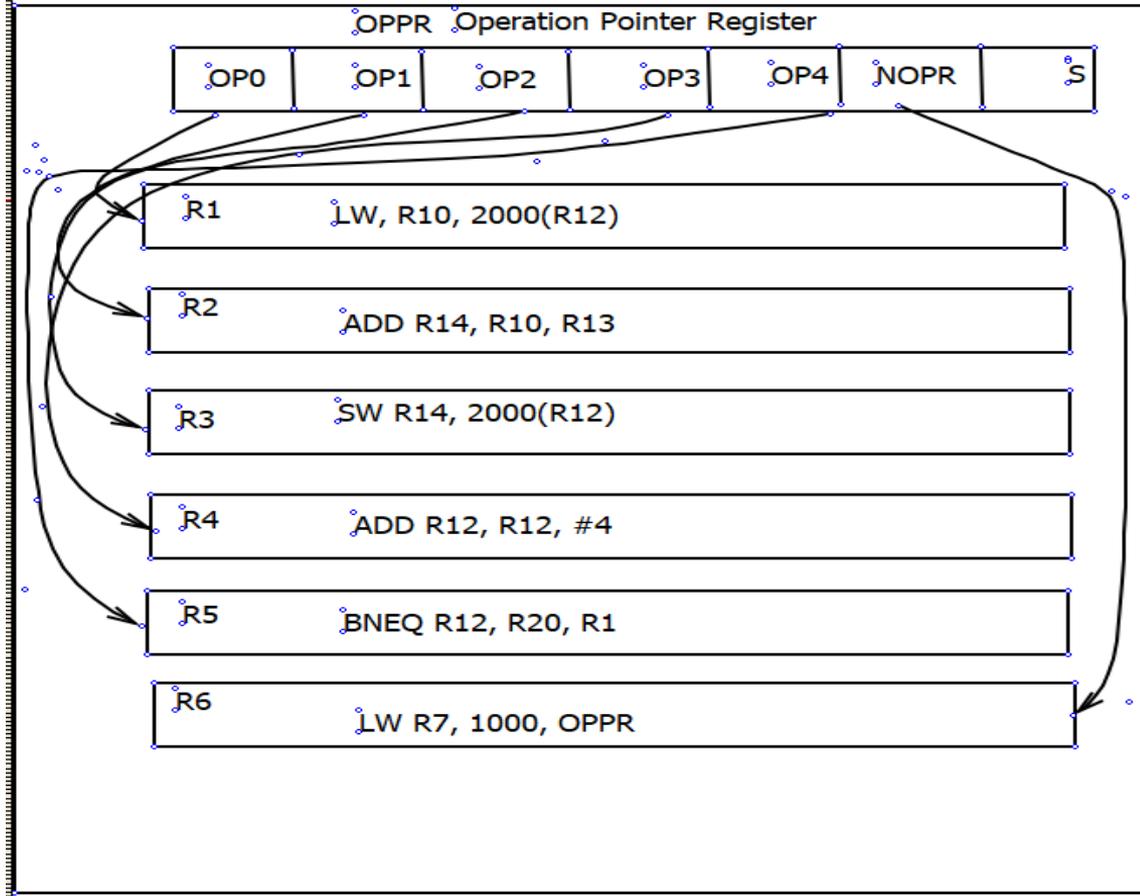
- It uses load to transfer operations and operands from memory to the register file
- It eliminates the use of a program counter that points only to one instruction at a time
- Single operation register pointer stores multiple pointers each points to one operation in another lower level registers
- The last field of register pointer points to the next operation register that loads the next block of code to be executed
- It executes multiple operations in parallel if the mode in operation pointer indicates a parallel operation

The Concept and Organization



Example of an application

- `for(i=0; i<N;i++){ A[i]=A[i]+B;}`.



Performance Evaluation

- `for(i=0; i<N;i++){A[i]=A[i]+B;}`. Register access=1 cycle, Memory access = 2 cycle

Conventional Architecture

Total time = $5N \times 1$ cycle for executing instructions + $2N \times 2$ cycles for data + $5N \times 2$ cycle for instructions memory = $19N$ cycles

Memory BW requirements = $(2N+5N) \times 4B / 19N = 1.5B/\text{cycle}$

No Instructions Computing

Total time = $5N \times 1$ cycle for executing instructions + 6×2 cycles to load the pointer and operations + $2N \times 2$ cycles for data load and store = $(9N + 12)$ cycles

Memory BW requirements = $2N \times 4B / (9N + 12) = .89 B/\text{cycle}$

Gain:

Performance gain = $19N / (9N + 12) = \text{about } 200\%$ and Memory BW reduction = $1.5B / .89 B = 168 \%$

Conclusions

- No Instructions Computing eliminates the use of program counter and reduces memory bandwidth requirements
- It is adaptable and supports SIMD and parallel operations Parallel operations from a single operation register pointer eliminates the overhead of synchronization
- Future work needs to implement it using suitable technology like FPGA and develop software tools to run actual applications

References

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2-Nagi Mekhiel, Take the 'No Instruction' road past memory straits, EE Times Aug 23, 1999